

# A Note on Weighted Rooted Trees

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## Abstract

Let  $T$  be a tree rooted at  $r$ . Two vertices of  $T$  are *related* if one is a descendant of the other; otherwise, they are *unrelated*. Two subsets  $A$  and  $B$  of  $V(T)$  are *unrelated* if, for any  $a \in A$  and  $b \in B$ ,  $a$  and  $b$  are unrelated. Let  $\omega$  be a nonnegative weight function defined on  $V(T)$  with  $\sum_{v \in V(T)} \omega(v) = 1$ . In this note, we prove that either there is an  $(r, u)$ -path  $P$  with  $\sum_{v \in V(P)} \omega(v) \geq \frac{1}{3}$  for some  $u \in V(T)$ , or there exist unrelated sets  $A, B \subseteq V(T)$  such that  $\sum_{a \in A} \omega(a) \geq \frac{1}{3}$  and  $\sum_{b \in B} \omega(b) \geq \frac{1}{3}$ . The bound  $\frac{1}{3}$  is tight. This answers a question posed in a very recent paper of Bonamy, Bousquet and Thomassé.

## 1 Introduction

Let  $T$  be a tree rooted at  $r$ . Let  $x \in V(T)$ . A *descendant* of  $x$  is any vertex  $y$  such that  $x \in V(P)$ , where  $P$  is the unique  $(r, y)$ -path in  $T$ . The *parent* of  $x$  is the vertex  $y$  such that  $y$  immediately precedes  $x$  on the unique  $(r, x)$ -path in  $T$ . Two vertices of  $T$  are *related* if one is a descendant of the other; otherwise, they are *unrelated*. Two subsets  $A$  and  $B$  of  $V(T)$  are *unrelated* if, for any  $a \in A$  and  $b \in B$ ,  $a$  and  $b$  are unrelated. Note that if  $A$  and  $B$  are unrelated, then  $A \cap B = \emptyset$ . Let  $G$  be a graph and let  $\omega$  be a nonnegative weight function defined on  $V(G)$ . For any  $A \subseteq V(G)$  and any subgraph  $H$  of  $G$ , define  $\omega(A) := \sum_{a \in A} \omega(a)$  and  $\omega(H) = \omega(V(H))$ . In their proof of the main result in [1], Bonamy, Bousquet and Thomassé made use of the following lemma.

**Lemma 1.1** *Let  $T$  be a tree rooted at  $r$  and let  $\omega$  be a nonnegative weight function defined on  $V(T)$  with  $\omega(T) = 1$ . Then there is an  $(r, u)$ -path  $P$  with  $\omega(P) \geq \frac{1}{4}$  for some  $u \in V(T)$ , or there exist unrelated sets  $A, B \subseteq V(T)$  such that  $\omega(A) \geq \frac{1}{4}$  and  $\omega(B) \geq \frac{1}{4}$ .*

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In the same paper, the authors believe that Lemma 1.1 holds for  $\frac{1}{3}$ . This problem has a Ramsey Theory flavor. In this note, we give an affirmative answer to their question and point out that the bound  $\frac{1}{3}$  is tight.

**Theorem 1.2** *Let  $T$  be a tree rooted at  $r$  and let  $\omega$  be a nonnegative weight function defined on  $V(T)$  with  $\omega(T) = 1$ . Then there is an  $(r, u)$ -path  $P$  with  $\omega(P) \geq \frac{1}{3}$  for some  $u \in V(T)$ , or there exist unrelated sets  $A, B \subseteq V(T)$  such that  $\omega(A) \geq \frac{1}{3}$  and  $\omega(B) \geq \frac{1}{3}$ . The bound  $\frac{1}{3}$  is tight.*

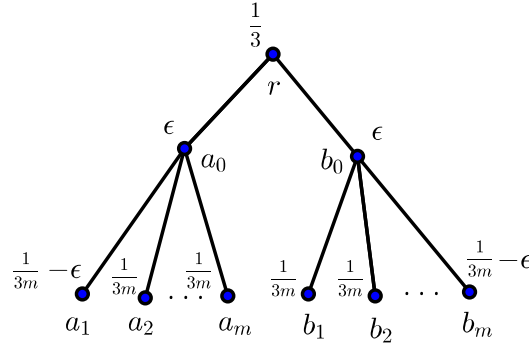


Figure 1: Rooted tree  $T$ .

To see why the bound  $\frac{1}{3}$  is best possible. Let  $m \geq 2$  be an integer and  $\epsilon \geq 0$  be a small number with  $\epsilon \leq \frac{1}{3m}$ . Let  $T$  be the weighted tree rooted at  $r$  as depicted in Figure 1. Note that  $\omega$  is a positive weight function on  $V(T)$  when  $\frac{1}{3m} > \epsilon > 0$ . Any path from the root  $r$  in  $T$  has weight between  $\frac{1}{3}$  and  $\frac{1}{3} + \frac{1}{3m} + \epsilon$ ; and  $T$  has one unique pair of unrelated sets  $A = \{a_0, a_1, a_2, \dots, a_m\}$  and  $B = \{b_0, b_1, b_2, \dots, b_m\}$  with  $\omega(A) = \omega(B) = \frac{1}{3}$ . The bound  $\frac{1}{3}$  is tight when  $m$  is large.

## 2 Proof of Theorem 1.2

Suppose  $T$  has no path from the root  $r$  with weight at least  $1/3$ . Then  $T$  is not a path. Let  $N_G(r) = \{v_1, v_2, \dots, v_s\}$  and  $T_1, T_2, \dots, T_s$  be connected components of  $T - r$ , where  $\omega(T_1) \leq \omega(T_2) \leq \dots \leq \omega(T_s)$ . We call each  $T_i$  a subtree of  $T$  rooted at  $v_i$  for  $1 \leq i \leq s$ . And  $T_1, \dots, T_s$  are also called *branches* of  $T$  at  $r$ . We next construct two unrelated sets  $A$  and  $B$  with desired weights according to the following algorithm:

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**Algorithm 1:** Building Sets  $A$  and  $B$ 

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**Data:** Vertex weighted tree  $T$  with root  $r$

**Result:** Unrelated sets  $A$  and  $B$  with desired weights

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1 Start at the root  $r$  with  $A = B = \emptyset$  and set  $C = \{T_1, T_2, \dots, T_s\}$ ;
2 while  $C \neq \emptyset$  do
3   for  $i = 1$  to  $s - 1$  do
4     Remove  $T_i$  from  $C$ . Add the vertices of  $T_i$  to  $A$  if  $\omega(A) \leq \omega(B)$ , and to  $B$ 
      otherwise;
5   end
6   If  $\emptyset \neq \bigcup_{i=1}^{s-1} V(T_i) \subseteq A$  (resp.  $\emptyset \neq \bigcup_{i=1}^{s-1} V(T_i) \subseteq B$ ), color the root  $r$  RED (resp.
      BLUE), otherwise color the root  $r$  GREEN ;
7   Set  $r$  to be the root of  $T_s$  and  $C$  be the set of connected components of  $T_s \setminus r$  with
      weights sorted in the nondecreasing order;
8 end
9 Call the last root  $r^*$ . If  $\omega(A) \leq \omega(B)$ , add  $r^*$  to  $A$  and color  $r^*$  RED, otherwise add
   $r^*$  to  $B$  and color  $r^*$  BLUE. Let  $y = r^*$ ,  $x$  be the parent of  $y$  and  $c$  be the color of  $y$ ;
10 while  $x$  is colored GREEN or  $c$  do
11   re-color  $x$  by the color  $c$  if  $x$  is colored GREEN and  $d_T(x) = 2$  ;
12   Set  $y$  to be  $x$ , and  $x$  be the parent of  $y$  ;
13 end
14 Let  $u = r^*$ ;
15 while  $u$  is adjacent to a vertex  $v \notin A \cup B$  with the same color as  $r^*$  do
16   Add  $v$  to  $A$  if both  $u$  and  $v$  are colored RED, and add  $v$  to  $B$  if both  $u$  and  $v$  are
      colored BLUE;
17   Set  $u$  to be  $v$ ;
18 end
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It can be easily checked that  $A$  and  $B$  constructed by the above algorithm are unrelated. Since  $T$  is not a path, both  $A$  and  $B$  are nonempty. Let  $u$  be the vertex in the last step of the algorithm that is added to  $A$  or  $B$ . According to the algorithm,  $u$  is colored RED or BLUE. Let  $M$  be the set of all colored vertices of  $T$ . Then the subgraph  $T[M]$  of  $T$  induced by  $M$  is the unique  $(r, r^*)$ -path, say  $P$ , where  $r^*$  is the last root as given in the algorithm. By the algorithm,  $T - A \cup B$  is the unique  $(r, u^*)$ -path, say  $P^*$ , of  $T$ , where  $u^*$  is the parent of  $u$  in  $T$ . Clearly,  $P^*$  is a subpath of  $P$ . Let  $N = V(P) - V(P^*)$ . Then  $r^* \in N$  and the vertices of  $N$  are all colored by the same color of the root  $r^*$ . One can see that if  $u$  is colored RED, then  $u \in N \subseteq A$  and the last set of vertices added to  $B$  are all uncolored. Similarly, if  $u$  is colored BLUE, then  $u \in N \subseteq B$  and the last set of vertices added to  $A$  are all uncolored. Since  $\omega(P^*) + \omega(A) + \omega(B) = 1$  and  $\omega(P^*) < \frac{1}{3}$ , we have

$$(1) \quad \omega(A) + \omega(B) > \frac{2}{3}.$$

We next show that  $\min\{\omega(A), \omega(B)\} \geq \frac{1}{3}$ .

Suppose that  $\omega(A) \leq \omega(B)$ . By (1),  $\omega(B) \geq \frac{1}{3}$ . Assume  $u \in B$ . Then  $u$  is colored BLUE and so  $r^*$  is also colored BLUE. Thus  $N \subseteq B$ . Since  $r^*$  is added to  $B$ , we have  $\omega(A) \geq \omega(B - N)$ . On the other hand,  $\omega(A) + \omega(B - N) = 1 - \omega(P) > \frac{2}{3}$ . Thus  $\omega(A) > \frac{1}{3}$ , as desired. So we may assume  $u \in A$ . Then  $u$  is colored RED and so  $N \subseteq A$ . Let  $D$  be the set of vertices that were last added to  $B$ . Then  $D$  contains only uncolored vertices of  $T$ . Thus  $D = V(Y)$ , where  $Y$  is a branch of some subtree  $T^*$  of  $T$ . Since  $D$  contains only uncolored vertices, by the algorithm,  $T^*$  has a branch  $X$  with  $\omega(Y) \leq \omega(X)$  and  $X \cap B = \emptyset$ . Let  $X^*$  be the set of all vertices that are added to  $A$  after the vertices in  $D$  were added to  $B$ . By the algorithm,  $X \subseteq X^*$ , and so  $\omega(X^*) \geq \omega(X) \geq \omega(Y)$ . Let  $\tilde{A} = A - X^*$  and  $\tilde{B} = B - Y$ . Since  $Y$  is added to  $B$ , we have  $\omega(\tilde{A}) \geq \omega(\tilde{B})$ . Note that  $\omega(A) = \omega(\tilde{A}) + \omega(X^*)$  and  $\omega(B) = \omega(\tilde{B}) + \omega(Y)$ . Thus  $\omega(B) \leq \omega(\tilde{A}) + \omega(X^*) = \omega(A)$ . By (1),  $\omega(A) \geq \frac{1}{3}$ . Hence  $\frac{1}{3} \leq \omega(A) \leq \omega(B)$ , as desired.

By a similar argument as above, one can show that  $\min\{\omega(A), \omega(B)\} \geq \frac{1}{3}$  for the case when  $\omega(B) \leq \omega(A)$ . This completes the proof of Theorem 1.2. ■

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## References

- [1] M. Bonamy, N. Bousquet, and S. Thomassé. The Erdős-Hajnal Conjecture for Long Holes and Antihole. <http://arxiv.org/pdf/1408.1964.pdf>, preprint, 2014